

## ERRATUM: DISTRIBUTED ANONYMOUS MOBILE ROBOTS: FORMATION OF GEOMETRIC PATTERNS\*

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**Abstract.** In this note we make a minor correction to a scheme for robots to broadcast their private information. All major results of the paper [I. Suzuki and M. Yamashita, *SIAM J. Comput.*, 28 (1999), pp. 1347–1363] hold with this correction.

**Key words.** anonymous robots, broadcast

**AMS subject classification.** 68Q99

**DOI.** 10.1137/050631562

**1. Correction.** Algorithms  $\psi_{f-point(2)}$  in section 3 and  $\psi_{getview}$  in section 4 of [1] use the following technique for all robots  $r_i$  to simultaneously “broadcast” to the other robots a privately chosen directed line  $\ell_i$  (e.g., the positive  $x$ -axis of its local coordinate system  $Z_i$ ). As outlined in the paragraph that follows the proof of Theorem 3.4, the basic idea is that each  $r_i$  moves repeatedly along  $\ell_i$  in a fixed direction each time it becomes active until, for each  $j \neq i$ , it has seen  $r_j$  at two or more distinct locations. Robot  $r_i$  can then figure out  $\ell_j$  based on two distinct locations that  $r_j$  has occupied. In an effort to ensure at the same time that  $r_j$  has also seen  $r_i$  at two or more distinct locations (so that it can figure out  $\ell_i$ ), we made an incorrect claim that if  $r_i$  has observed  $r_j$  at three or more distinct locations, then  $r_j$  has observed  $r_i$  at two or more distinct locations. This claim must be replaced by the following.

**PROPOSITION 1.** *For any integer  $m \geq 1$ , if  $r_i$  has seen  $r_j$  at  $2m$  distinct locations in the time interval  $[0, t]$ , then  $r_j$  has seen  $r_i$  at  $m$  or more distinct locations in  $[0, t]$ .*

*Proof.* Suppose  $r_i$  becomes active at times  $t_1$  and  $t_2$ ,  $t_1 < t_2$ , and observes  $r_j$  at two distinct locations. Since  $r_j$  occupies distinct locations at  $t_1$  and  $t_2$ ,  $r_j$  must become active in  $[t_1, t_2]$  and observe  $r_i$  at a location on the line segment  $\overline{p_1p_2}$ , where  $p_1$  and  $p_2$  are the locations of  $r_i$  on line  $\ell_i$  at  $t_1$  and  $t_2$ , respectively. This means that  $r_j$  observes  $r_i$  at a distinct location each time  $r_i$  observes  $r_j$  at two distinct locations, since  $r_i$  moves along  $\ell_i$  in a fixed direction each time it becomes active.  $\square$

Therefore, to ensure that  $r_j$  has seen  $r_i$  at two distinct locations,  $r_i$  must continue to move along  $\ell_i$  until it has seen  $r_j$  at four or more distinct locations. However, if we allow  $r_i$  to simply stop moving as soon as it has observed  $r_j$  at four or more distinct locations, then  $r_j$  may not be able to observe  $r_i$  at four or more distinct locations. This means that  $r_j$  may never finish the broadcast.

Fortunately, broadcasting a line is usually a preliminary step that precedes a main task. In the case of  $\psi_{f-point(2)}$ , the ultimate goal of  $r_i$  and  $r_j$  is to move to the midpoint  $p$  of their initial positions. (Because of the way  $\ell_i$  and  $\ell_j$  are chosen in  $\psi_{f-point(2)}$ , both  $r_i$  and  $r_j$  can compute  $p$  from  $\ell_i$  and  $\ell_j$ .) Note that when  $r_i$  has seen  $r_j$  at four or more distinct locations,  $r_i$  knows not only  $\ell_j$ , but also that  $r_j$  knows  $\ell_i$  (because by Proposition 1  $r_j$  has seen  $r_i$  at two or more distinct locations). Thus  $r_i$

\*Received by the editors May 13, 2005; accepted for publication (in revised form) February 22, 2006; published electronically June 19, 2006.

<http://www.siam.org/journals/sicomp/36-1/63156.html>

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can now safely quit broadcasting  $\ell_i$  and move to (or toward)  $p$ , regardless of whether  $r_j$  has seen  $r_i$  at four or more distinct locations. By moving to a point not on  $\ell_i$ ,  $r_i$  effectively “announces” to  $r_j$  that it now knows  $\ell_j$ . Robot  $r_j$  eventually observes  $r_i$  at a location not on  $\ell_i$  and learns that it can quit broadcasting  $\ell_j$  as well and proceed to  $p$ .

Based on the above discussion,  $r_i$  can use the following scheme to broadcast  $\ell_i$  to all other robots:

1.  $r_i$  moves along  $\ell_i$  in a fixed direction each time it becomes active until, for each  $j \neq i$ , either
  - (a) it has seen  $r_j$  at four or more distinct locations, or
  - (b) it observes that  $r_j$  is at a location not on  $\ell_j$ . (Note that  $r_i$  knows  $\ell_j$  by the time this case occurs.)
2. Then it moves to a point not on  $\ell_i$ .

Algorithm  $\psi_{f-point(2)}$ , the proof of Theorem 3.3, and the paragraph that follows the proof of Theorem 3.4 should be revised accordingly.

Algorithm  $\psi_{getview}$  must be modified using a similar idea. In  $\psi_{getview}$ , each robot  $r_i$ , initially located at the origin  $o_i$  of its local coordinate system  $Z_i$ , broadcasts three lines: its  $x$ -axis,  $y$ -axis, and line  $L_i$  through  $o_i$  in direction  $f(d_i)$ , where  $d_i$  is the minimum distance between any two initial positions of the robots, and for  $x > 0$ ,  $f(x) = (1 - 1/2^x) \times 90^\circ$  is a monotonically increasing function with range  $(0^\circ, 90^\circ)$ . (This function replaces  $f(x) = (1 - 1/2^x) \times 360^\circ$  used in the paper. Both  $f(d_i)$  and  $d_i$  are measured in terms of  $Z_i$ .) Note that the orientations of the three lines are all distinct.

Each robot  $r_i$  first broadcasts its  $x$ -axis by moving along it in the positive direction. When  $r_i$  knows that all other robots know  $r_i$ ’s  $x$ -axis and its orientation, i.e., for each  $j \neq i$ , either

- (a)  $r_i$  has seen  $r_j$  at four or more distinct locations, or
- (b)  $r_i$  observes that  $r_j$  has changed the direction of its motion ( $r_i$  knows the  $x$ -axis of  $r_j$  by the time this occurs),

it returns straight to  $o_i$  and starts broadcasting its  $y$ -axis by moving along it in the positive direction (thereby changing its direction of motion and announcing to others that it now knows the  $x$ -axes of all other robots). Eventually all robots finish broadcasting their  $x$ -axes and start broadcasting their  $y$ -axes.  $r_i$  ends the broadcast of its  $y$ -axis when it knows that all other robots know its  $y$ -axis and its orientation, returns straight to  $o_i$ , and starts broadcasting line  $L_i$  by moving in direction  $f(d_i)$  (thereby changing its direction of motion again, announcing to others that it knows the  $y$ -axes of all other robots). Eventually all robots  $r_j$  start broadcasting their lines  $L_j$ . Once again,  $r_i$  terminates this broadcast when it knows that all other robots know  $L_i$ , and returns to its initial position  $o_i$ . By returning to  $o_i$ ,  $r_i$  announces to others that it has finished the broadcast of  $L_i$ .

#### REFERENCE

[1] I. SUZUKI AND M. YAMASHITA, *Distributed anonymous mobile robots: Formation of geometric patterns*, SIAM J. Comput., 28 (1999), pp. 1347–1363.